Distribution of Patterns in Pairwise Sequence Alignments

S. Wolfsheimer and G. Nuel

PARIS DESCARTES MAP5 – Mathématiques Appliquées á Paris 5, University Paris Descartes, France

Sequence Alignment [1]

Detection of homological relationships in DNA- or protein-sequences. Applications: Search tools for molecular

databases, e.g. BLAST



- Given: Pair of sequences $\mathbf{a} = a_1 \dots a_M \in \Sigma^M$ $\mathbf{b} = b_1 \dots b_N \in \Sigma^N.$
- amino acids: $\Sigma = \{A, C, D, E \dots\}$ nucleotides: $\Sigma = \{A, G, C, T\}$
- Alignment \mathcal{A} : set of pairings $(i_1, j_1), \ldots, (i_n, j_n)$ with $i_k < i_{k+1}$, $j_k < j_{k+1}$
- Gaps: insertions and deletions of subsequences



global gapped alignment

Alignment score:

$$\mathsf{S}(\mathcal{A};\mathbf{a},\mathbf{b}) = \sum_{\mathsf{pairs}} \sigma_{a_i,b_j} - \sum_{\mathsf{gaps}} g(l_{\mathsf{gap}})$$

• score matrix: $\sigma_{a,b} = \log \frac{P_{a,b}}{f_a f_b}$ gap penalty: $g(l) = \alpha + \beta(l-1)$

Score based alignment: Similarity between a and b measured by optimal alignment score S_0 .

- $S_0 = \max_{\mathcal{A}} \mathsf{S}(\mathcal{A}; \mathbf{a}, \mathbf{b})$
- $\mathcal{A}_0 = \operatorname{argmax}_4 S(\mathcal{A}; \mathbf{a}, \mathbf{b})$

Dynamic Programming

Global alignment: Needleman-Wunsch algorithm

> • Here: linear gap costs $g(l_{gap}) =$ δl_{gap}

Affine gap-costs straight forward.

• A_{ij} : optimal alignment of subproblem $a_1, ..., a_i, b_1, ..., b_j$

$$A_{i,j} = \max \begin{cases} A_{i-1,j-1} + \sigma_{a_i,b_j} \\ A_{i-1,j} - \delta \\ A_{i,j-1} - \delta \end{cases}$$

So = Amn

• $\mathcal{O}(M \cdot N)$ time complexity

Alignment biases [3]

Close to gaps often many competitive alignments decrease the accuracy of the score based alignment. Typical effects are:

• Gap wander: a gap shifted by a few positions.

A TTCTATGCCGGCAGTG / TTCTGCG----GCG

 Gap attraction: two close gaps merge into a single gap.

В ТССАGСАТGСТGGССС // ТСТА--АТGС--GССС T C C A G C A T G C T G G C C C T C T A A T G C - - - - G C C C

• Gap annihilation: two gaps of opposite signature (insertion / deletion) cancel each other.

D CCT -- ATGCGTATGCATGCC CCTATGCGTATGCATGCC

Probabilistic alignment: more quantitative description. Here: distributions of typical patterns that might lead to weakly reliable alignment segments.

Probabilistic alignment [2]

- $P_T(\mathcal{A})$: distribution of global alignments of a. b.
- · Canonical ensemble of alignments
 $$\begin{split} P_T(\mathcal{A};\mathbf{a},\mathbf{b}) &= \frac{1}{Z_T} \exp\left[\frac{1}{T}S(\mathcal{A};\mathbf{a},\mathbf{b})\right] \\ \left(\text{Boltzmann distribution with "tem-} \right. \end{split}$$
 perature" T = 1)
- Partition function: $Z_T = \sum_{\mathcal{A}} \exp\left[S(\mathcal{A})/T\right]$
- Partition function calculation: $D_{i,j} \to Z_{i,j}, \max \to \sum, \text{ and } + \to \times$
- Forward algorithm: $Z_{i,j}$: sum over all alignments of $a_1 \dots a_i$ and $b_1 \dots b_i$

$$Z_{i,j} = \underbrace{Z_{i-1,j-1}e^{\sigma_{a_i,b_j}/T}}_{\text{match/mismatch}} + \underbrace{(Z_{i-1,j} + Z_{i,j-1})e^{-\delta/2}}_{\text{match/mismatch}}$$

$$Z_T = Z_{M,I}$$

• Backward algorithm: $Z'_{i,i}$: sum over all alignments of $a_{i+1} \ldots a_M$ and $b_{i+1} \ldots b_N$ given that $(i, j) \in \mathcal{A}$.

• Posterior probabilities:

$$p_{ij} := P[(i, j) \in \mathcal{A}] = \frac{1}{Z_T} Z_{i,j} Z'_{i,j}$$

 $P\begin{pmatrix}a_i\\\end{pmatrix} = 1 - \sum p_{ij}$

$$Y \in \mathcal{A}.$$

or probabilities:
 $P[(i,j) \in \mathcal{A}] = \frac{1}{Z_T} Z_{i,j} Z'_{i,j}$

$$P\left(\begin{array}{c}a_i\\-\end{array}\right) = 1 - \sum_j p_{ij}$$

GTGTTAGT-----AGTGAAGTAGGTCGATG-GTGTAGC---AGC |||||||| GTGTTAGTAGTTCCAGTGATTTAGGATG-TGTGTGTAGCTTTAGC

Number of gaps

Distribution of the number of gaps (see [4] for a general HMM setup). $a_{i-l} \cdots a_{i-1} a_i \quad a_{i+1} \cdots a_{i+l'}$ Insertion b_i b_{i+l}

Deletion a_i $- a_{i+1}$

 $b_{j-l} \cdots b_{j-1} b_j \quad b_{j+1} \cdots b_{j+l'}$ Insertions: auxiliary matrices $Z_{i,j}^{(I,k)}$ and $I_{i,j}^{(I,k)}$ for $k = 1, 2, \ldots$, sum over alignments, k insertions seen so far. $\tau(I,k)$ $\ldots \ldots a_i$

$$\begin{split} I_{i,j}^{(i,j)'} &: \text{sum over} \qquad \dots \quad b_j \quad - \quad - \\ Z_{i,j}^{(I,k)} &= \left(I_{i-1,j-1}^{(I,k)} + Z_{i-1,j-1}^{(I,k)} \right) e^{\sigma_{a_i,b_j}/T} \\ &+ \left(Z_{i,j-1}^{(I,k)} + I_{i,j-1}^{(I,k)} \right) e^{-\delta/T} \\ I_{i,j}^{(I,k)} &= \left(Z_{i-1,j}^{(I,k-1)} + I_{i-1,j}^{(I,k)} \right) e^{-\delta/T} \\ P(k) &= \left(Z_{M,N}^{(I,k)} + I_{M,N}^{(I,k)} \right) / Z_{M,N} \\ P\left(\begin{array}{c} a_i \\ - \end{array}, \quad kth \text{ insertion} \right) &= \sum_j \frac{I_{i,j}^{(I,k)} \cdot I_{i,j+1}'}{Z_{M,N}} \\ k = 1 \\ k = 2 \\ \dots \\ k = 6 \\ \text{TGTTAGT} \\ \text{$$

Gap Annihilation Pattern

Detecting	candidate	s for	gap	anni-
nilation.	Search	for t	the p	attern:
Insertion	··· –	a_{i-k+1}	$1 \cdots a_i$	a_{i+1}
	$\cdots b_{j-k}$	b_{j-k+1}	$1 \cdots b_j$	-

Deletion
$$\cdots a_{i-k} a_{i-k+1} \cdots a_i - \cdots - \underbrace{b_{j-k+1} \cdots b_j}_{j-k+1} b_{j+1}$$

Decode alternative alignment segments:



Bibliography

- [1] Durbin et.al., Biological sequence analysis (1998)
- [2] Miyazawa., Prot. Eng. 8 (1995)
- [3] Lunter et.al. , Genome Research. 18 (2008) [4] Aston and Martin, Ann. Appl. Stat.1 (2007)